

# Bad semidefinite programs: they all look the same

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## A pair of Semidefinite Programs (SDP)

$$\begin{array}{ll} \sup_x c^T x & \inf_Y B \bullet Y \\ \text{s.t. } \sum_{i=1}^m x_i A_i \preceq B & Y \succeq 0 \\ & A_i \bullet Y = c_i \quad (i = 1, \dots, m). \end{array}$$

Here

- $A_i, B$  are symmetric matrices,  $c, x \in \mathbb{R}^m$ .
- $A \preceq B$  means that  $B - A$  is symmetric positive semidefinite (psd).
- $A \bullet B = \sum_{i,j} a_{ij} b_{ij}$ .

## Conic LPs and SDPs

- Common framework for LP and SDP: both  $\mathbb{R}_+^n$  and **psd matrices** are closed convex **cones**.
- A set  $C$  is a **cone**, if  $x \in C, \lambda \geq 0 \Rightarrow \lambda x \in C$ .
- Linear objective, conic constraint both in LP and SDP, and many other interesting problems, notably SOCPs.

## Why is SDP important: applications in

- 0–1 Integer programming.
- Approx algorithms
- Chemical engineering
- Chemistry
- Coding theory
- Control theory
- Combinatorial opt
- Discrete geometry
- Eigenvalue optimization
- Facility planning
- Finance
- Geometric optimization
- Global optimization
- Graph visualization
- Inventory theory
- Machine learning
- Matrix analysis
- PDEs
- Probability theory
- Robust optimization
- Signal processing
- Statistics
- Structural optimization

# SDP theory and applications

- **Nice duality theory:** see later
- **Applications:** see textbooks by
  - **Boyd-Vandenberghe**
  - **Ben-Tal-Nemirovskii**
- **Algebraic geometry:**
  - **Nie-Sturmfels 2010**
  - **von Bothmer-Ranestad 2009**
  - **Gouveia, Parrilo, Thomas, 2010**
  - **Book by Blekherman, et al, 2013**
- **Polynomial optimization:**
  - **Lasserre 2000 –**
  - **Parrilo 2000 –**
  - **Nie 2000 –**
  - **Helton-Vinnikov 2003**

# SDP duality

The primal-dual pair of SDPs:

$$\sup_x c^T x$$

$$s.t. \sum_{i=1}^m x_i A_i \preceq B$$

$$\inf_Y B \bullet Y$$

$$Y \succeq 0$$

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**But:** in SDP, unlike in LP **pathological phenomena** occur: nonattainment, positive gaps.

This is bad, since we would like a certificate of optimality.

## Pathology # 1: nonattainment in dual

Primal:

$$\begin{array}{ll} \sup 2x_1 & \\ \text{s.t. } x_1 \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \preceq \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} & \Leftrightarrow \sup 2x_1 \\ & \text{s.t. } \begin{pmatrix} 1 & -x_1 \\ -x_1 & 0 \end{pmatrix} \preceq 0 \end{array}$$

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Here  $\inf = 0$ , but not attained: Any  $y_{11} > 0$ ,  $y_{22} = 1/y_{11}$  is feasible, but  $y_{11} = 0$  is not.

## Pathology # 2: positive duality gap

Primal:

$\sup x_2$

$$s.t. \quad x_1 \begin{pmatrix} 1 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix} + x_2 \begin{pmatrix} 0 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \end{pmatrix} \preceq \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 0 \end{pmatrix}$$

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Only feasible  $x_2$  is  $x_2 = 0$ .

Dual value is **1**, and it is attained.



# Terminology

## Definition:

- The system  $P_{SD} = \{ x \mid \sum_{i=1}^m x_i A_i \preceq B \}$  is **well-behaved**, if for all  $c$  such that

$$\sup\{ c^T x \mid x \in P_{SD} \} \text{ is finite,}$$

the dual program has the same value, and it attains.

- **Badly behaved**, otherwise.
- We would like to understand well/badly behaved systems.

## Some literature

- Conic LPs may be badly behaved when  $K$  is not polyhedral.
- **Borwein-Wolkowicz 1981** Facial reduction: theoretical construction of well behaved system
- **Ramana 1995** Extended dual for SDP
- **Ramana, Tunçel, Wolkowicz, 1997** Facial reduction implies correctness of extended dual
- **Klep, Schweighofer, 2013** Related duals based on algebraic geometry.
- **Waki, Muramatsu, 2013; Pataki 2013:** Simpler facial reduction algorithms.
- **P 2007** Closedness of linear image of a closed, convex cone

## Motivation

The systems

$$x_1 \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \preceq \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$$

and

$$x_1 \begin{pmatrix} 1 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix} + x_2 \begin{pmatrix} 0 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \end{pmatrix} \preceq \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 0 \end{pmatrix}$$

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are both badly behaved.

Curious similarity:

- “Hanging off” diagonals;
- if we delete 2nd row and 2nd column in all matrices in the second system, and delete the first matrix, we get back the first system.

## Why all bad SDPs look the same

- Semidefinite system:

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- W.l.o.g. the max (rank) slack is

$$Z = \begin{pmatrix} I_r & 0 \\ 0 & 0 \end{pmatrix}.$$

Then  $(P_{SD})$  badly behaved  $\Leftrightarrow \exists V$  a lin. combination of the  $A_i$  and  $B$  as

$$V = \begin{pmatrix} \overbrace{V_{11}}^r & V_{12} \\ V_{12}^T & V_{22} \end{pmatrix}, \text{ where } V_{22} \succeq 0, \mathbf{R}(V_{12}^T) \not\subseteq \mathbf{R}(V_{22}).$$

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- Ex:  $x_1 \overbrace{\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}}^V \preceq \overbrace{\begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}}^Z$

## What is missing?

- Matrices  $Z, V$  prove that  $(P_{SD})$  is badly behaved.
- But: this is not yet a poly time, or easy to verify proof of bad behavior

## Reformulations of

$$(PSD) \sum_{i=1}^m x_i A_i \preceq B$$

are obtained by a sequence of:

- Rotate all matrices by  $T = \begin{pmatrix} I_r & 0 \\ 0 & M \end{pmatrix}$ ,  $M$  orthogonal.
- $B \leftarrow B + \sum_{i=1}^m \mu_i A_i$
- $A_i \leftarrow \sum_{j=1}^m \lambda_j A_j$  where  $\lambda_i \neq 0$

Reformulations preserve well/badly behaved status; preserve max slack; provide an equivalence relation

Theorem:  $(P_{SD})$  is badly behaved  $\Leftrightarrow$  it has a reformulation:

$$(P_{SD,bad}) \quad \sum_{i=1}^k x_i \begin{pmatrix} F_i & 0 \\ 0 & 0 \end{pmatrix} + \sum_{i=k+1}^m x_i \begin{pmatrix} F_i & G_i \\ G_i^T & H_i \end{pmatrix} \preceq \begin{pmatrix} I_r & 0 \\ 0 & 0 \end{pmatrix} = Z,$$

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Proof that  $(P_{SD,bad})$  is badly behaved:

$$x \text{ feas. with slack } S \Rightarrow \text{supp}(S) \subseteq \text{supp}(Z)$$

$$\Rightarrow x_{k+1} = \dots = x_m = 0 \Rightarrow \text{sup } -x_m = 0$$

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Proof that  $(P_{SD,bad})$  is badly behaved:

$$\begin{aligned} Y \succeq 0, Y \bullet Z = 0 &\Rightarrow Y = \begin{pmatrix} 0 & 0 \\ 0 & Y_{22} \end{pmatrix} \\ &\Rightarrow Y \bullet \begin{pmatrix} F_m & G_m \\ G_m^T & H_m \end{pmatrix} \geq 0 \\ &\Rightarrow \text{no dual soln with value } 0 \end{aligned}$$



## Example: before reformulation

$$\begin{aligned} & \begin{matrix} x_1 \\ +x_2 \\ +x_3 \\ +x_4 \end{matrix} \begin{pmatrix} 54 & 46 & 50 & 4 \\ 46 & -38 & 87 & -106 \\ 50 & 87 & -60 & 296 \\ 4 & -106 & 296 & -368 \end{pmatrix} + \begin{pmatrix} 110 & 91 & 105 & -6 \\ 91 & -72 & 171 & -210 \\ 105 & 171 & -72 & 528 \\ -6 & -210 & 528 & -672 \end{pmatrix} + \begin{pmatrix} 42 & 35 & 40 & 0 \\ 35 & -28 & 67 & -82 \\ 40 & 67 & -36 & 216 \\ 0 & -82 & 216 & -272 \end{pmatrix} \\ & \qquad \qquad \qquad + \begin{pmatrix} 36 & 30 & 35 & -2 \\ 30 & -24 & 57 & -70 \\ 35 & 57 & -24 & 176 \\ -2 & -70 & 176 & -224 \end{pmatrix} \approx \begin{pmatrix} 389 & 323 & 370 & -12 \\ 323 & -257 & 610 & -748 \\ 370 & 610 & -288 & 1920 \\ -12 & -748 & 1920 & -2432 \end{pmatrix} \end{aligned}$$

Hard to tell if well or badly behaved

## Example: after reformulation

$$\begin{aligned}
 & x_1 \left( \begin{array}{cc|cc} 0 & 1 & 0 & 0 \\ 1 & -2 & 0 & 0 \\ \hline 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) + x_2 \left( \begin{array}{cc|cc} 2 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ \hline 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{array} \right) + x_3 \left( \begin{array}{cc|cc} 0 & 0 & 2 & 1 \\ 0 & 0 & 3 & -1 \\ \hline 2 & 3 & 0 & 2 \\ 1 & -1 & 2 & 0 \end{array} \right) \\
 & + x_4 \left( \begin{array}{cc|cc} 0 & 0 & 3 & -1 \\ 0 & 0 & 2 & -1 \\ \hline 3 & 2 & 4 & 0 \\ -1 & -1 & 0 & 0 \end{array} \right) \quad | \quad \mathcal{L} \quad \left( \begin{array}{cc|cc} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ \hline 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{array} \right)
 \end{aligned}$$

As before:  $x_3 = x_4 = 0 \Rightarrow \sup -x_4 = 0$

But: no dual solution with value 0

Theorem:  $(P_{SD})$  is well behaved  $\Leftrightarrow$  it has a reformulation:

$$(P_{SD,good}) \quad \sum_{i=1}^k x_i \begin{pmatrix} F_i & 0 \\ 0 & 0 \end{pmatrix} + \sum_{i=k+1}^m x_i \begin{pmatrix} F_i & G_i \\ G_i^T & H_i \end{pmatrix} \preceq \begin{pmatrix} I_r & 0 \\ 0 & 0 \end{pmatrix} = Z,$$

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- $(P_{SD})$  well behaved  $\Rightarrow$  for all  $c$  with a finite obj. value  $\exists$  optimal

$$Y = \begin{pmatrix} \overbrace{Y_{11}}^r & 0 \\ 0 & Y_{22} \end{pmatrix}$$

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- **Corollary:** we can generate **all** well behaved semidefinite systems: choose in sequence  $H_i, G_i, F_i$ . Then do reformulation.

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  - **Corollary:** we can generate **all** linear maps under which the image of the psd cone is closed.

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- **Corollary:** we can generate **all** well behaved semidefinite systems: choose in sequence  $H_i, G_i, F_i$ . Then do reformulation.
  - **Corollary:** we can generate **all** linear maps under which the image of the psd cone is closed.
  - **Proof:**  $\{(A_i \bullet Y)_{i=1}^m \mid Y \succeq 0\}$  is closed  $\Leftrightarrow \sum_{i=1}^m x_i A_i \preceq 0$  is well behaved.

## Broader framework: Well- and badly behaved conic LPs

- Conic linear system, with  $K$  closed, convex cone:

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- Ex: LP, SDP, SOCP, ...

## Well- and badly behaved conic LPs

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- Conic LP:

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- Ex: LP, SDP, SOCP, ...

- $(P)$  is **well-behaved**, if

$$\sup(P_c) = \min(D_c) \forall c$$

where  $(D_c)$  is dual program. **Badly behaved** if not well behaved.

## Well- and badly behaved conic LPs

- Conic linear system, with  $K$  closed, convex cone:

$$(P) \quad Ax \leq_K b \quad (\Leftrightarrow b - Ax \in K)$$

- Conic LP:

$$(P_c) \quad \sup \{ \langle c, x \rangle \mid x \text{ feasible in } (P) \}$$

- Ex: LP, SDP, SOCP, ...

- $(P)$  is **well-behaved**, if

$$\sup(P_c) = \min(D_c) \forall c$$

where  $(D_c)$  is dual program. **Badly behaved** if not well behaved.

- **Known:**

$K$  polyhedral  $\Rightarrow (P)$  is well-behaved.

$(P)$  Slater, i.e.,  $\exists x : b - Ax \in \text{ri } K \Rightarrow (P)$  is well-behaved.

Given  $K$  closed, convex cone

- dual cone:  $K^* = \{ y \mid \langle x, y \rangle \geq 0 \forall x \in K \}$
- $K$  is nice if  $K^* + F^\perp$  is closed for all  $F$  faces of  $K$ .

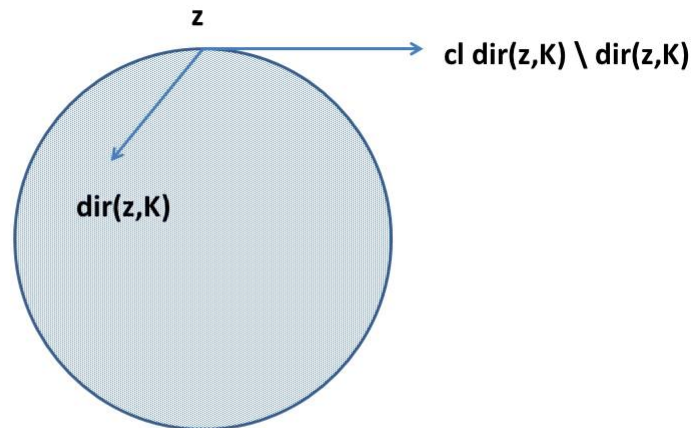


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- set of feasible directions at  $z \in K$  – maybe not closed:  
 $\text{dir}(z, K) = \{ y \mid \exists \epsilon > 0 \text{ s.t. } z + \epsilon y \in K \}$



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- $\Rightarrow$  is true, even if  $K$  is not nice
- $K$  polyhedral, or  $(P)$  Slater ( $z \in \text{ri}K$ )  $\Rightarrow \text{dir}(z, K)$  closed.

## The difference set

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$$V \in \text{cl dir}(Z, PSD) \setminus \text{dir}(Z, PSD)$$

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$$V = \begin{pmatrix} \overbrace{V_{11}}^r & V_{12} \\ V_{12}^T & V_{22} \end{pmatrix}, \text{ where } V_{22} \succeq 0, \text{R}(V_{12}^T) \not\subseteq \text{R}(V_{22}).$$

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- We recover characterization of badly behaved semidefinite systems.

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  - **Generating** all linear maps under which the image of the psd cone is closed.
- More generally: conditions for well and badly behaved nature of a conic linear system
- Exact characterization when  **$K$**  is **nice**.
- Latest version of paper is on Optimization Online.

Thank you!